

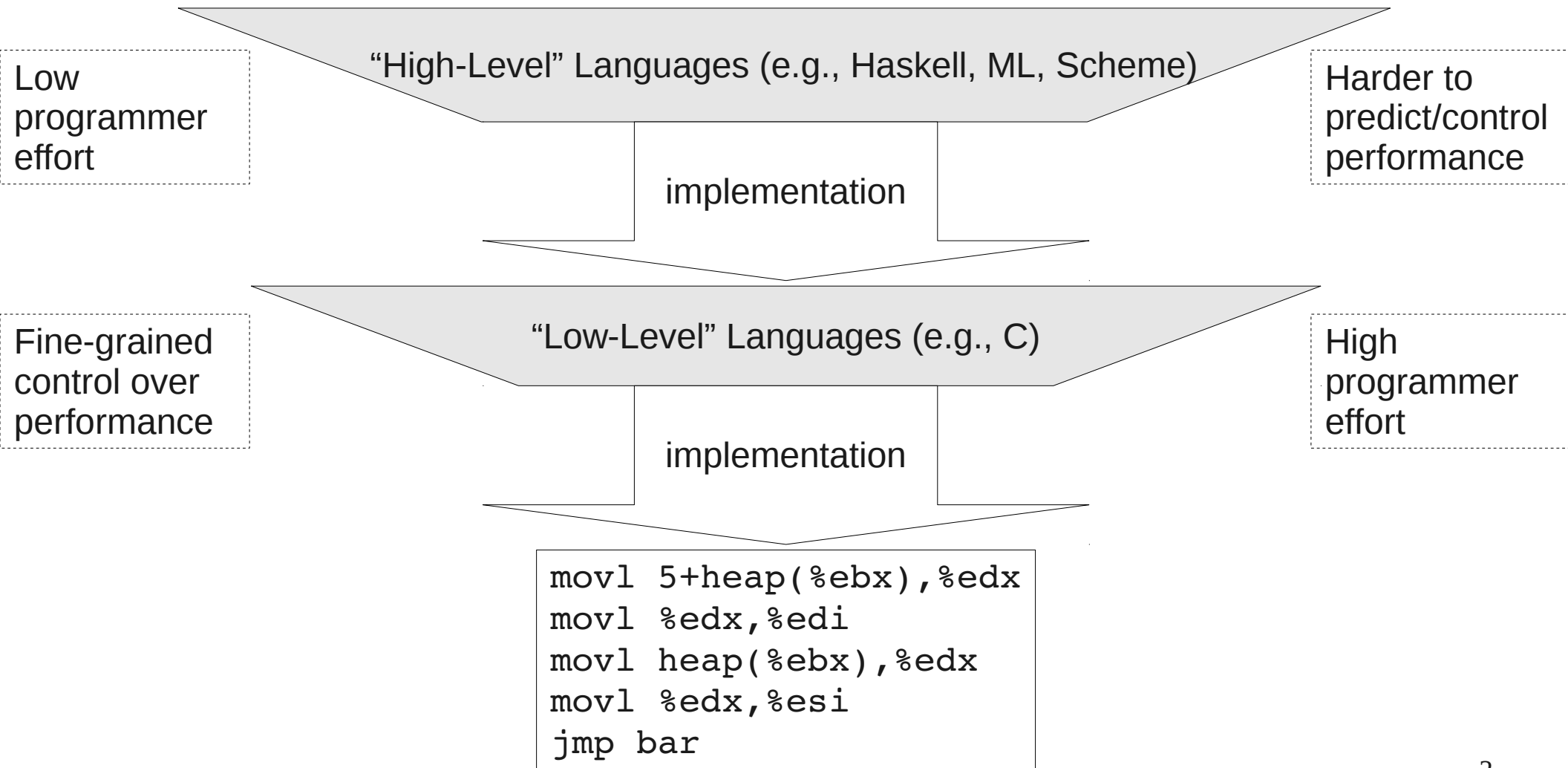
# The Bedrock Structured Programming System

Combining Generative  
Metaprogramming and Hoare  
Logic in an Extensible  
Program Verifier

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ICFP 2013  
September 27, 2013

# In the beginning, there was assembly language....



# The high cost of abstraction

$$1 + 1$$

vs.

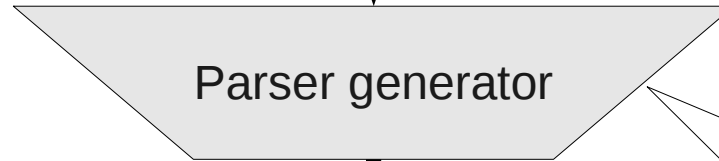
$$\begin{array}{l} \mathbf{f} \ \mathbf{x} \ \mathbf{y} = \mathbf{x} + \mathbf{y} \\ \mathbf{f} \ 1 \ 1 \end{array}$$

Is there a performance cost to functional abstraction?  
higher-order functions?  
modules?  
laziness?  
garbage collection?  
exceptions?

# Having your cake & eating it, too: code generators

```
A ::= "(" B ")" | C
B ::= "foo" | "bar"
C ::= "baz" | C A
```

We pay in *performance* for this abstraction gap only at **compile time!**



Conventional implementations with string munging are awfully hard to get right....

Steaming pile of C code

assembly

# Enter embedded domain-specific languages!

```
A ::= "( B )" | C
B ::= "foo" | "bar"
C ::= "baz" | C A
```

Parser generator

Steaming pile of C code

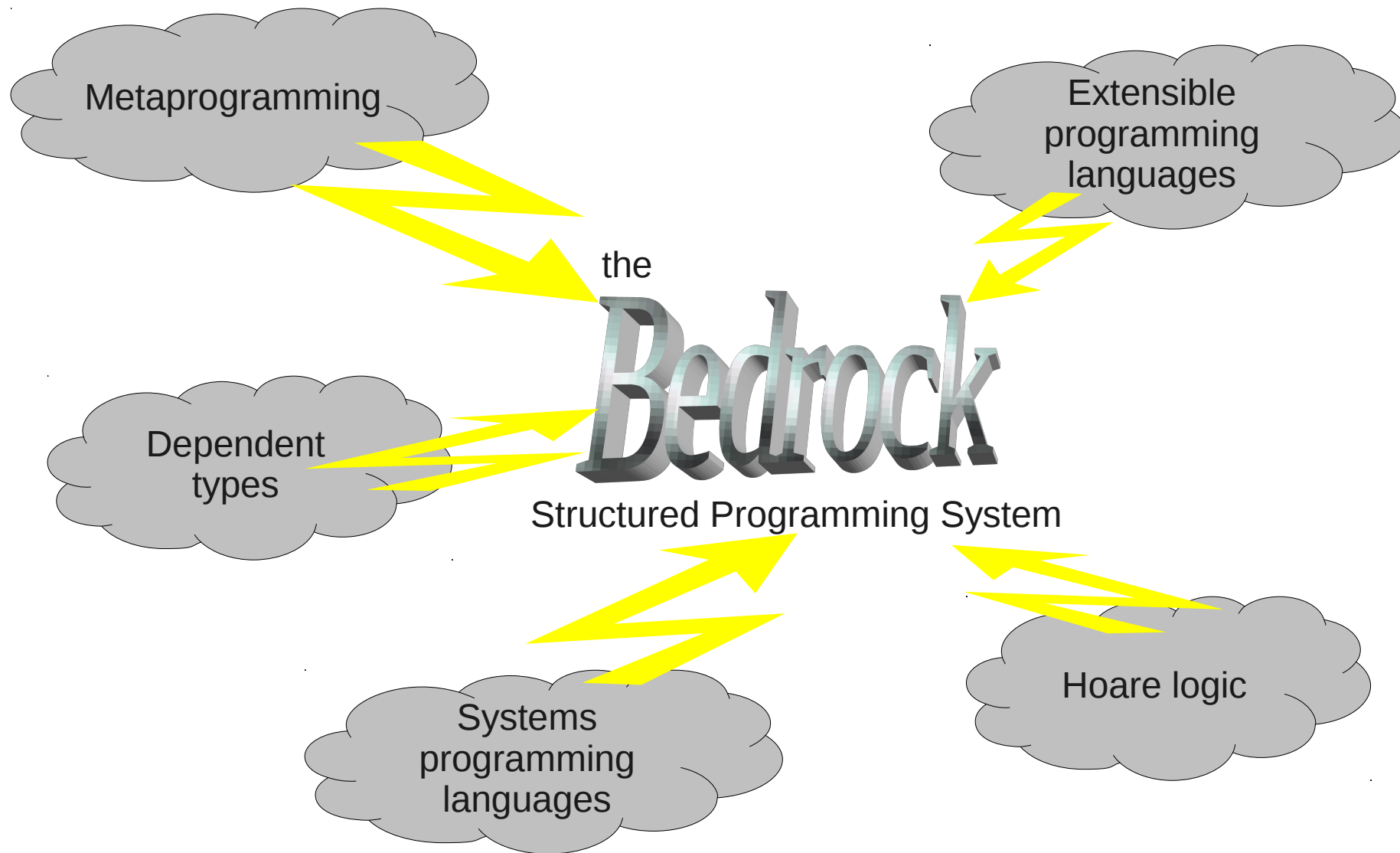
assembly

The type system of the **metalanguage** (e.g., Haskell) helps us guarantee that the translation always generates reasonable code in the **object language** (e.g., C)!

- ✓ Basic syntactic well-formedness
- ✓ Variable binding
- ✓ Type checking
- ✓ **Functional correctness...?**

Haskell term of type  
Grammar  $\rightarrow$  CProgram

# What this talk is probably about



Results in this talk use other parts of Bedrock contributed by Gregory Malecha, Thomas Braibant, Patrick Hulin, and Edward Z. Yang!

# The big picture



Inside Coq

The **Bedrock IL**

Very low-level program representation

assembly

Functional program, drawing on libraries of dependently typed combinators

Ingredients for verifying that program against specifications in higher-order logic

**Essential goal:**

From these ingredients, we can verify the program *without knowing how the code generator works!*

# The Bedrock IL

$W ::= (* \text{ width-32 bitvectors } *)$   
 $L ::= (* \text{ program code block labels } *)$

$\text{Reg} ::= \text{Sp} \mid \text{Rp} \mid \text{Rv}$   
 $\text{Loc} ::= \text{Reg} \mid W \mid \text{Reg} + W$   
 $\text{Lvalue} ::= \text{Reg} \mid [\text{Loc}]_{32} \mid [\text{Loc}]_8$   
 $\text{Rvalue} ::= \text{Lvalue} \mid W \mid L$   
 $\text{Binop} ::= + \mid - \mid *$   
 $\text{Test} ::= = \mid != \mid < \mid <=$

$\text{Instr} ::= \text{Lvalue} := \text{Rvalue} \mid \text{Lvalue} := \text{Rvalue} \text{ Binop} \text{Rvalue}$

$\text{Jump} ::= \text{goto} \text{Rvalue} \mid \text{if} \text{Rvalue} \text{Test} \text{Rvalue} \text{ then goto } L \text{ else goto } L$

$\text{Block} ::= \text{Instr}^*; \text{Jump}$   
 $\text{Spec} ::= (* \text{ assertion language of XCAP } *)$   
 $\text{Module} ::= (L: \{\text{Spec}\} \text{Block})^*$

Why not LLVM or a similar IR?

*Answer:* Builds in a host of features:

- ◆ Types
- ◆ Variables
- ◆ Functions

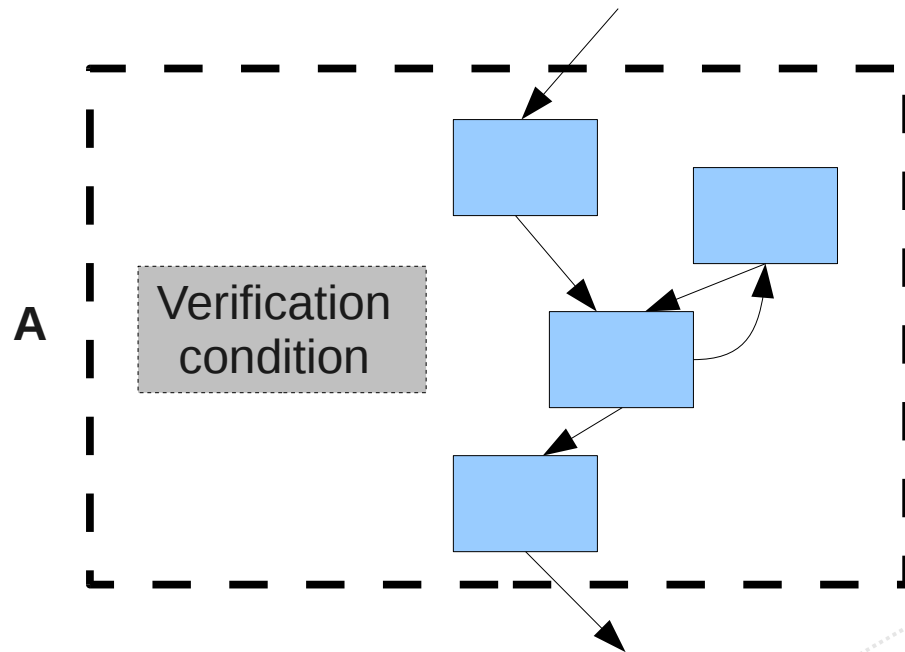
We will implement all of these as **libraries** in Bedrock!

A simple language makes it easier to prove *foundational program correctness* theorems in Coq.

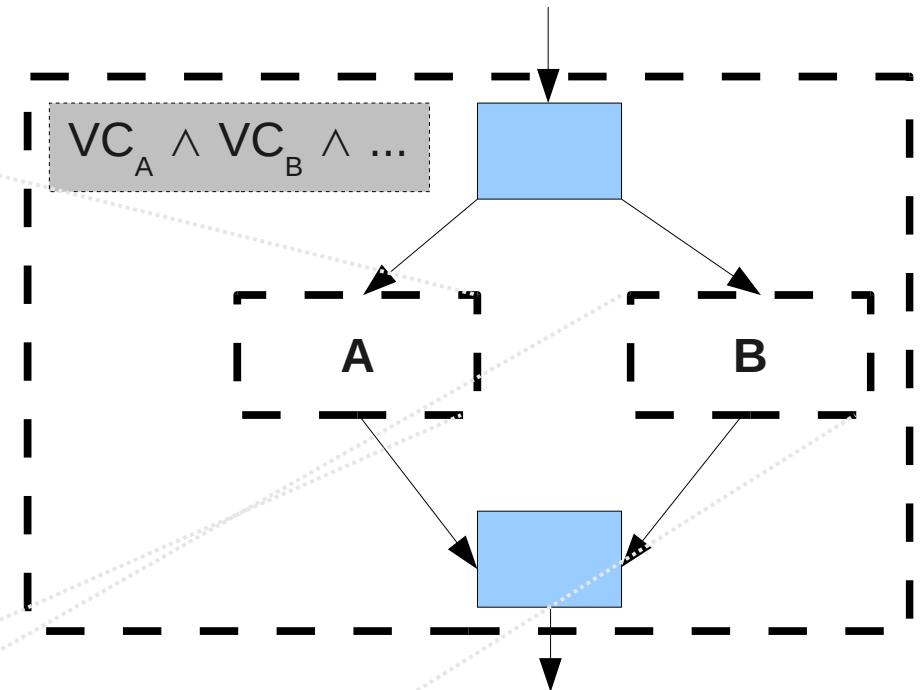


# One-slide sketch of the formal details

An extensible C-like language based on *macros* that produce *chunks* encapsulating control-flow graphs:



if..then..else combinator (functional program):



Details omitted here:

Plumbing of connecting CFGs  
Predicate transformers  
Formal connection to Hoare logic

# Bedrock version of linked list length

Specification

```
Definition lengthS : spec := SPEC("x") reserving 1
  Al ls,
  PRE[V] sll ls (V "x")
  POST[R] [| R = length ls |] * sll ls (V "x").
```

Program

```
bfunction "length"("x", "n") [lengthS]
  "n" <- 0;; Loop invariant
  [Al ls,
    PRE[V] sll ls (V "x")
    POST[R] [| R = V "n" ^+ length ls |] * sll ls (V "x")]
  While ("x" <> 0) {
    "n" <- "n" + 1;;
    "x" <-* "x" + 4
  };;
  Return "n"
end.
```

This is all Coq code, taking advantage of Coq's extensible parser!

Proof

```
Theorem sllMOk : moduleOk sllM.
Proof.
  vcgen; abstract (sep hints; finish).
Qed.
```

# Pattern matching for network protocols

```
"pos" <- 0;;  
Match "req" Size "len" Position "pos" {  
  Case (0 ++ "x")  
    Return "x"  
  end;;  
  Case (1 ++ "x" ++ "y")  
    Return "x" + "y"  
  end  
} Default {  
  Fail  
}
```

# Declarative querying of arrays

```
"acc" <- 0;;
```

Fancy macro-specific loop invariant form

```
[After prefix Approaching all
```

```
PRE[V] [| V "acc" = countNonzero prefix |]
```

```
POST[R] [| R = countNonzero all | ]
```

```
For "index" Holding "value" in "arr" Size "len"
```

```
Where (Value <> 0) {
```

```
  "acc" <- "acc" + 1
```

```
};;
```

```
Return "acc"
```

Loop has filter condition that the macro analyzes syntactically to decide on optimizations.

# Build toolchain

Coq program of “chunk” type



Apply Coq function from Bedrock library

Coq program of “list of basic blocks” type



Coq program extraction

OCaml program of “list of basic blocks” type



OCaml execution (with side effects)

.s assembly file

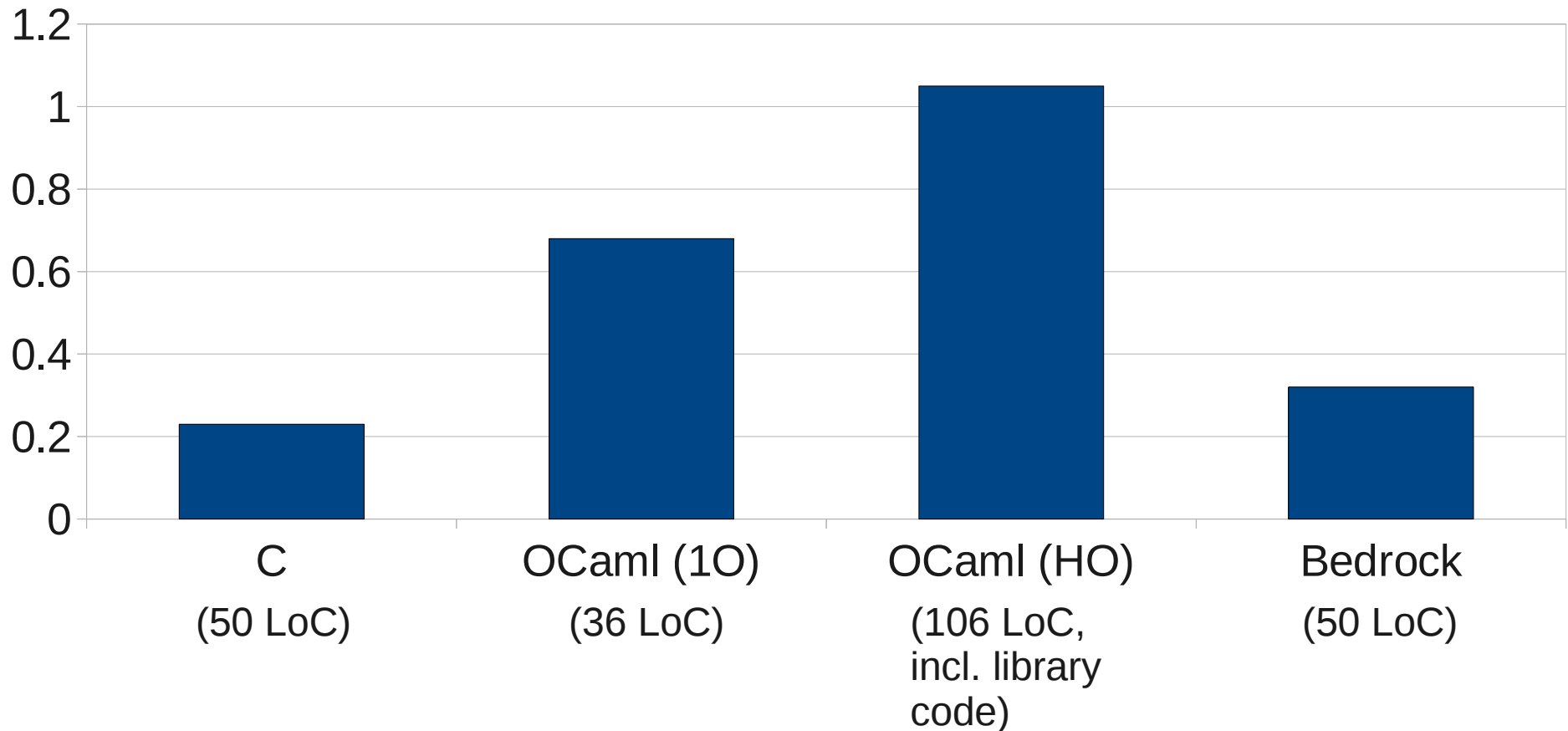


Normal GNU build tools

ELF binary

# Running time comparison on a database-inspired benchmark

All programs parse and execute the same set of 200 random queries over a random array of length 100,000.



# Bedrock on the web

`http://plv.csail.mit.edu/bedrock/`